



**Politecnico di Bari**  
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**C3LAB**  
Control of Computing  
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Systems Lab

# A Hybrid Model of the Akamai Adaptive Streaming Control System

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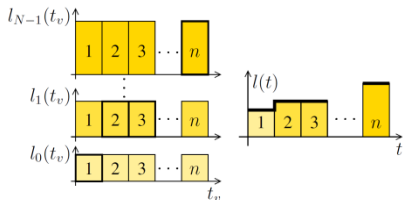
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1. Introduction
2. The Akamai Control System
3. Model and Properties
4. Experimental Validation
5. Conclusion

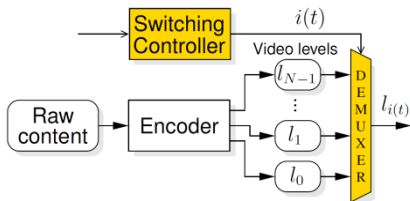


- In video streaming a client streams the video over an HTTP connection from the server that stores it;
- Video streaming is becoming the largest fraction of the Internet traffic;
- With the stream-switching approach the video bitrate can be throttled on-the-fly to match the time-varying available bandwidth;
- At the client a buffer is employed to absorb instantaneous bandwidth variations;

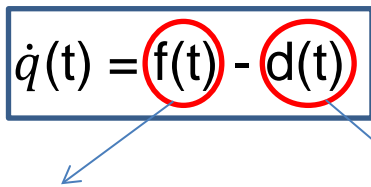
# THE STREAM-SWITCHING APPROACH



- The raw video is encoded in N different qualities (video levels or bitrates) which are stored at the server;
- Each level is divided into segments of fixed duration;
- At every segment download the Switching Controller selects the level of the next segment;



In a generic video streaming control system the playout buffer length, i.e. the total duration of video stored in the playout buffer, can be modelled as:

$$\dot{q}(t) = f(t) - d(t)$$


filling rate:

$$f(t) = \frac{r(t)}{l(t)}$$

$r(t)$  is the received rate,  $l(t)$  is the selected video bitrate

draining rate:

$$d(t) = \begin{cases} 1, & \text{playing} \\ 0, & \text{paused} \end{cases}$$



## Motivation: why are we considering the Akamai stream-switching control system?

- it is a **leading** CDN operator whose platform is employed by several streaming platforms, including Livestream;
- it presents some **interesting and perhaps unique control features**;

... however, the code is **not open source**, thus this work is based on a previous work of system identification!

(see De Cicco, L. and Mascolo, S., "*An adaptive video streaming control system: Modeling, validation, and performance evaluation*". IEEE/ACM Transactions on Networking)...



The Akamai control system consists of *two controllers*:

1. the *stream-switching controller*, which selects the video level;
2. the *playout buffer length controller*, which aims at avoiding that the buffer gets empty (buffering events) by throttling the sending rate;

Their behavior changes according to the logical phase in which the control system is. There are *three phases*:

1. the *Buffering phase*: entered to quickly fill the queue at the start or after a buffering event;
2. the *Normal phase* (periodically triggered): the sending rate is throttled to steer the buffer length to a target length;
3. the *Greedy phase* (periodically triggered): the sending rate is set to a much higher value than in the *Normal* phase to probe and estimate the available bandwidth;

The *stream-switching controller*.

- given the available bandwidth, selects the optimal video level  $l_{opt}(t)$  such that:

$$\begin{array}{l}
 l_{opt}(t) = \underset{l \in L}{\operatorname{argmax}} l \\
 \text{s. t. } \hat{b}(t) > (1 + S_f)l \\
 S_f > 0
 \end{array}$$

estimated bandwidth ←

→ Safety margin

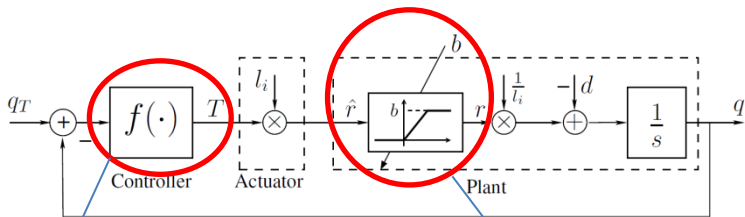
- is event-based:
  - level switch-up is triggered when  $\hat{b}(t) > (1 + S_f)l_{i+1}$
  - level switch-down is triggered when  $q(t) < q_d$

lower threshold



# THE AKAMAI PLOUT BUFFER LENGTH CONTROLLER

The *plout buffer length controller* throttles the sending rate to steer the buffer length to the target  $q_T$ :



$T(t)$  is equal to:

1.  $T_B > 1$  (*Buffering phase*);
2.  $\max(1 + \frac{q_T - q(t)}{q_T}, T_m)$  (*Normal phase*);
3.  $T_M > 2$  (*Greedy phase*);

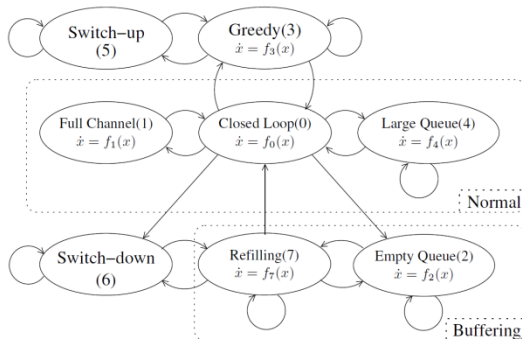
the saturation block models the bottleneck link of bandwidth  $b(t)$



- A piecewise constant bandwidth input function has been employed. It allows us to analyze any practical traffic scenario with a bottleneck link;
- The communication forward and backward delays from the client to the server and the actuation time-delay have been neglected;
- The safety factor  $S_f$ , the queue threshold  $q_d$  and the queue target  $q_T$  have been assumed to be constant (in the reality they are time-varying);

# THE HYBRID AUTOMATON

Due to the **state-dependent** and **event-triggered dynamics** and the **discontinuous elements** the model has the form of the following hybrid automaton:



The state  $x = [i, r, q, \tau, \tau_1, \tau_2]^T$ :

- $i$  (the video level index);
- $r$  (the sending rate);
- $q$  (the queue length);
- $\tau$  (the phases timer),
- $\tau_1$  (time-varying duration of the Normal phase timer);
- $\tau_2$  (time-varying duration of the Greedy phase timer);

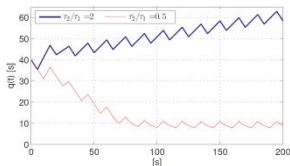
# PROPERTIES

Proposition 1: a necessary condition to permit a switch-up between two adjacent levels  $l_i$  and  $l_{i+1}$  is that  $\frac{l_{i+1}-l_i}{l_i} \leq \frac{T_M}{S_f}$ .

Remark: this condition ensures the **reachability of all the video levels**.

Proposition 2: a sufficient condition for the boundedness of the playout buffer length  $q(t)$  is that  $\frac{\tau_2}{\tau_1} \leq \frac{1-T_m}{T_M-1}$ .

Remark: **large buffering** is a waste of resources from the network point of view. It is prevented when this condition is satisfied.



Proposition 3: let us assume that an actuation time-delay  $\tau_a$  occurs when a switch-down event is triggered and that  $B \geq l_0$ . A sufficient condition to avoid buffering is that  $q_d > (1 - \frac{l_0}{l_i})\tau_a$ .

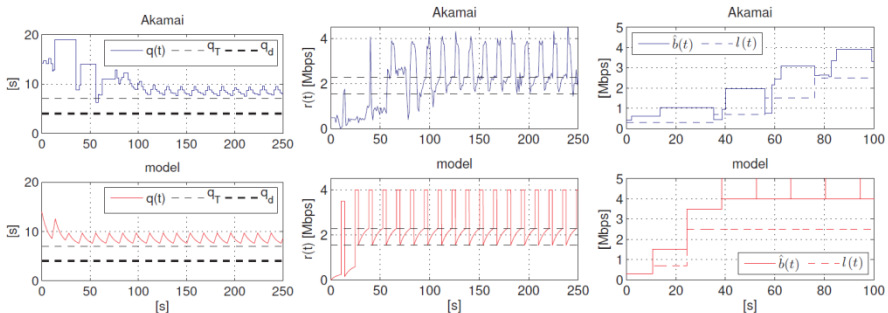
Remark: thanks to this property the threshold  $q_d$  can be tuned to ensure **robustness against actuation time-delays**.

## THE MODEL VALIDATION PROCEDURE



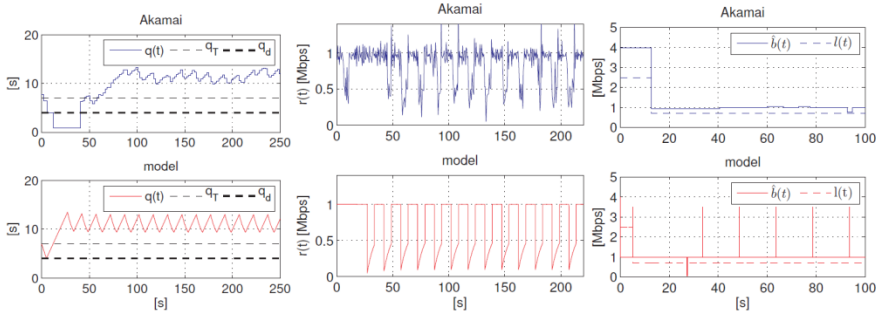
- the dynamics of the variables of the simulated model have been compared with the ones obtained through Internet experiments;
- the model has been simulated with the Matlab *Hybrid Equations (HyEq) Toolbox*;
- the experimental results have been obtained by playing the video “Elephant’s Dream”, served by the Akamai server, on a Linux PC equipped with the traffic shaper **tc** to change the link capacity in real time;
- abrupt step-like bandwidth increases and decreases have been considered to validate, respectively, the switch-up and the switch-down case;

# THE SWITCH-UP CASE



- The model is quite accurate;
- The succession of the two phases is confirmed;
- The difference in the transient is due to the neglected actuation delay, which is present in the real system;

# THE SWITCH-DOWN CASE



- The model is quite accurate;
- The alternating dynamics due to phases are shown;
- Here, too, the difference in the transient is due to the neglected actuation delay;

## CONCLUSIONS

- A model of the Akamai stream-switching control system has been proposed;
- The model is in the form of a hybrid automaton;
- The model has been validated by comparing simulations and experimental results;
- We have provided insights on the parameters tuning by means of some key properties;



# Questions?